Introduction to Natural Language Processing I [Statistické metody zpracování přirozených jazyků I] (NPFL067)

http://ufal.mff.cuni.cz/courses/npfl067

Word Classes: Programming Tips & Tricks

The Algorithm (review)

- Define merge(r,k,l) = (r',C') such that
 - C' = C $\{k,l\} \cup \{m \text{ (a new class)}\}$
 - r'(w) = r(w) except for k,l member words for which it is m.
- 1. Start with each word in its own class (C = V), r = id.
- 2. Merge two classes k,l into one, m, such that

 $(k,l) = \operatorname{argmax}_{k,l} I_{\operatorname{merge}(r,k,l)}(D,E).$

- 3. Set new (r,C) = merge(r,k,l).
- 4. Repeat 2 and 3 until |C| reaches a predetermined size.

Complexity Issues

- Still too complex:
 - |V| iterations of the steps 2 and 3.
 - $|V|^2$ steps to maximize $\operatorname{argmax}_{k,l}$ (selecting k,l freely from |C|, which is in the order of $|V|^2$)
 - |V|² steps to compute I(D,E) (sum within sum, all classes, also: includes log)
 - \Rightarrow total: $|V|^5$
 - i.e., for |V| = 100, about 10^{10} steps; ~ several hours!
 - but $|V| \sim 50,000$ or more

Trick #1: Recomputing The MI the Smart Way: Subtracting...

• Bigram count table:



- Test-merging c_2 and c_4 : recompute only rows/cols 2 & 4:
 - subtract column/row (2 & 4) from the MI sum (intersect.!)
 - add sums of merged counts (row & column)

...and Adding

• Add the merged counts:



Trick #2: Precompute the Counts-to-be-Subtracted

- Summing loop goes through i,j
- ...but the single row/column sums do not depend on the (resulting sums after the) merge
- \Rightarrow can be precomputed
 - only 2k logs to compute at each algorithm iteration, instead of k²
- Then for each "merge-to-be" compute only add-on sums, plus "intersection adjustment"

Formulas for Tricks #1 and #2

Let's have <u>k</u> classes at a certain iteration. Define:
 q_k(l,r) = p_k(l,r) log(p_k(l,r) / (p_{kl}(l) p_{kr}(r)))
 now the same, but using counts:

 $q_k(l,r) = c_k(l,r)/N \log(N c_k(l,r)/(c_{kl}(l) c_{kr}(r)))$

• Define further (row+column <u>i</u> sum): precomputed $s_k(a) = \sum_{l=1..k} q_k(l,a) + \sum_{r=1..k} q_k(a,r) - q_k(a,a)$ • Then, the subtraction part of Trick #1 amounts to $sub_k(a,b) = s_k(a) + s_k(b) - q_k(a,b) - q_k(b,a)$ remaining intersect. adj.

Formulas - cont.

• After-merge add-on:

 $add_{k}(a,b) = \sum_{l=1..k, l \neq a, b} q_{k}(l,a+b) + \sum_{r=1..k, r \neq a, b} q_{k}(a+b,r) + q_{k}(a+b,a+b)$

- What is it <u>a+b</u>? Answer: the <u>new (merged) class</u>.
- Hint: use the definition of q_k as a "macro", and then $p_k(a+b,r) = p_k(a,r) + p_k(b,r)$ (same for other sums, equivalent)
- The above sums cannot be precomputed
- After-merge Mutual Information (I_k is the "old" MI, kept from previous iteration of the algorithm):

 $I_k(a,b)$ (MI after merge of cl. a,b) = I_k - sub_k(a,b) + add_k(a,b)

Trick #3: Ignore Zero Counts

- Many bigrams are 0
 - (see the paper: Canadian Hansards, < .1 % of bigrams are non-zero)
- Create linked lists of non-zero counts in columns and rows (similar effect: use perl's hashes)
- Update links after merge (after step 3)

Trick #4: Use Updated Loss of MI

- We are now down to |V|⁴: |V| merges, each merge takes |V|² "test-merges", each test-merge involves order-of-|V| operations (add_k(i,j) term, foil #8)
- <u>Observation</u>: many numbers (s_k, q_k) needed to compute the mutual information loss due to a merge of i+j *do not change*: namely, those which are not in the vicinity of neither i nor j.
- <u>*Idea*</u>: keep the MI loss matrix for all pairs of classes, and (after a merge) update only those cells which have been influenced by the merge.

Formulas for Trick #4 (s_{k-1}, L_{k-1})

- Keep a matrix of "losses" $L_k(d,e)$.¹
- Init: $L_k(d,e) = sub_k(d,e) add_k(d,e)$ [then $I_k(d,e) = I_k L_k(d,e)$]
- Suppose a,b are now the two classes merged into a:
- Update (k-1: index used for the <u>*next*</u> iteration; $i, j \neq a, b$):

 $\begin{aligned} &-s_{k-1}(i) = s_k(i) - q_k(i,a) - q_k(a,i) - q_k(i,b) - q_k(b,i) + q_{k-1}(a,i) + q_{k-1}(i,a) \\ &- {}^2L_{k-1}(i,j) = L_k(i,j) - s_k(i) + s_{k-1}(i) - s_k(j) + s_{k-1}(j) + q_k(i+j,a) + q_k(a,i+j) + q_k(i+j,b) + q_k(b,i+j) - \end{aligned}$

- $q_{k-1}(i+j,a)$ - $q_{k-1}(a,i+j)$ [NB: may substitute even for s_k , s_{k-1}]

NB ¹ L_k is symmetrical L_k(d,e) = L_k(e,d) (q_k is something different!) ²The update formula L_{k-1}(l,m) is wrong in the Brown et. al paper

Completing Trick #4

- $s_{k-1}(a)$ must be computed using the "Init" sum.
- $L_{k-1}(a,i) = L_{k-1}(i,a)$ must be computed in a similar way, for all $i \neq a,b$.
- $s_{k-1}(b)$, $L_{k-1}(b,i)$, $L_{k-1}(i,b)$ are not needed anymore (keep track of such data, i.e. mark every class already merged into some other class and do not use it anymore).
- Keep track of the minimal loss during the $L_k(i,j)$ update process (so that the next merge to be taken is obvious immediately after finishing the update step).

Efficient Implementation

- Data Structures: (N # of bigrams in data [fixed])
 - Hist(k) history of merges
 - Hist(k) = (a,b) merged when the remaining number of classes was k
 - $c_k(i,j)$ bigram class counts [updated]
 - $c_{kl}(i), c_{kr}(i)$ unigram (marginal) counts [updated]
 - $-L_k(a,b)$ table of losses; upper-right trianlge [updated]
 - $s_k(a)$ "subtraction" subterms [optionally updated]
 - $q_k(i,j)$ subterms involving a log [opt. updated]
 - The optionally updated data structures will give linear improvement only in the subsequent steps, but at least $s_k(i)$ is necessary in the initialization phase (1st iteration)

Implementation: the Initialization Phase

- 1 Read data in, init counts $c_k(l,r)$; then $\forall l,r,a,b$; a < b:
- 2 Init unigram counts:

 $c_{kl}(l) = \sum_{r=1..k} c_k(l,r), \quad c_{kr}(r) = \sum_{l=1..k} c_k(l,r)$

- complicated? remember, must take care of start & end of data!

- 3 Init $q_k(l,r)$: use the 2nd formula (count-based) on foil 7, $q_k(l,r) = c_k(l,r)/N \log(N c_k(l,r)/(c_{kl}(l) c_{kr}(r)))$
- 4 Init $s_k(a) = \sum_{l=1..k} q_k(l,a) + \sum_{r=1..k} q_k(a,r) q_k(a,a)$

• 5 Init
$$L_k(a,b) = s_k(a) + s_k(b) - q_k(a,b) - q_k(b,a) - q_k(a+b,a+b) +$$

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$$\sum_{l=1..k, l \neq a, b} q_k(l, a+b)$$
 - $\sum_{r=1..k, r \neq a, b} q_k(a+b, r)$

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Implementation: Select & Update

- 6 Select the best pair (a,b) to merge into <u>a</u> (watch the candidates when computing L_k(a,b)); save to Hist(k)
- 7 Optionally, update q_k(i,j) for all i,j ≠ b, get q_{k-1}(i,j)
 remember those q_k(i,j) values needed for the updates below
- 8 Optionally, update s_k(i) for all i ≠ b, to get s_{k-1}(i)
 again, remember the s_k(i) values for the "loss table" update
- 9 Update the loss table, $L_k(i,j)$, to $L_{k-1}(i,j)$, using the tabulated q_k , q_{k-1} , s_k and s_{k-1} values, or compute the needed $q_k(i,j)$ and $q_{k-1}(i,j)$ values dynamically from the counts: $c_k(i+j,b) = c_k(i,b) + c_k(j,b)$; $c_{k-1}(a,i) = c_k(a+b,i)$

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Towards the Next Iteration

- 10 During the L_k(i,j) update, keep track of the minimal loss of MI, and the two classes which caused it.
- 11 Remember such best merge in Hist(k).
- 12 Get rid of all s_k , q_k , L_k values.
- 13 Set k = k -1; stop if k == 1.
- 14 Start the next iteration
 - either by the optional updates (steps 7 and 8), or
 - directly updating $L_k(i,j)$ again (step 9).

Moving Words Around

- Improving Mutual Information
 - take a word from one class, move it to another (i.e., two classes change: the moved-from and the moved-to), compute $I_{new}(D,E)$; keep change permanent if $I_{new}(D,E) > I(D,E)$

keep moving words until no move improves I(D,E)

- Do it at every iteration, or at every <u>m</u> iterations
- Use similar "smart" methods as for merging

Using the Hierarchy

- Natural Form of Classes
 - follows from the sequence of merges:



evaluation assessment analysis understanding opinion

Numbering the Classes (within the Hierarchy)

• Binary branching

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• Assign 0/1 to the left/right branch at every node:

